Lecture 25: Atomic Locks COSC 273: Parallel and Distributed Computing Spring 2023

Announcements

Homework 03 is finalized

- no new questions
- due next Friday

Today

• More Lock Implementations

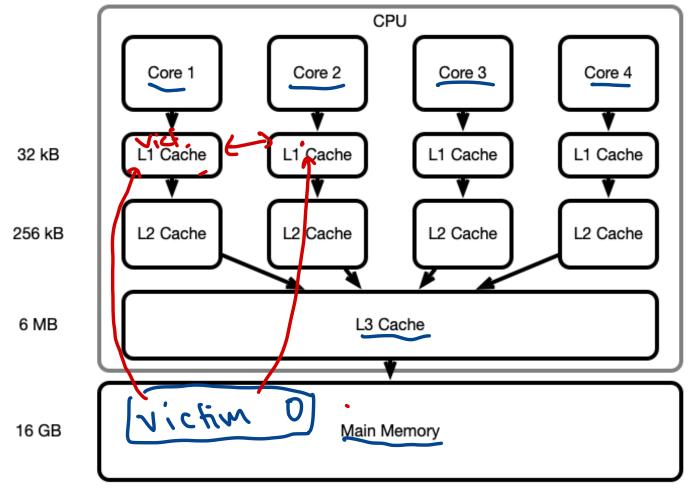
Last Time:

- Peterson lock implementation
 - peterson-lock.zip
- disappointment
 - it didn't achieve mutual exclusion!

Peterson Lock Code

Memory Consistency!

" cache Coherence 16



volatile Variables

Java can make variables visible between threads:

- use volatile keyword
- individual read/write operations to volatile are atomic
 Drawbacks:
- volatile variables are less efficient
- *only* single read/write operations are atomic
 - e.g. count++ not atomic
- only primitive datatypes are visible
 - if volatile SomeClass..., only the reference is treated as volatile

Making Variables Volatile

- In PetersonLock
 - flag: an *array* (object) can't be volatile
 replace with boolean flag0, flag1
 - victim <</pre>
- In LockedCounter
 - count

Fixing Implementation

• peteson-lock.zip

Finally!!!

What have we done?

- 1. Proven correctness of a lock
 - idealized model of computation
 - atomic read/write operations
- 2. Implemented lock
 - used Java to resemble idealized model
- 3. Used lock
 - saw expected behavior

Theory and practice converge!

Peterson: Good and Bad

The Good:

1. It works!

2. It only uses read/write operations!

The Bad:

- 1. It only works with two threads!
- 2. Ugly implementation
 - need a separate PetersonThread to assign IDs

Question. How could we lock more simply?

Better Tech!

Use more advanced Atomic Objects!

Introducing the AtomicBoolean class:

- var ab = new AtomicBoolean(boolean value) make an AtomicBoolean with initial value value
- [ab.get() return the current value]
- ab.getAndSet(boolean newValue) atomically set the value to newValue and return the old value
- ab.compareAndSet(boolean expected, boolean new) atomically update to new if previous value was expected and return whether or not the value was updated
 if (value == expected)

if (value == expected) Value = new return true else return false

A Simpler Lock?

Question. How could we use AtomicBooleans to design a simpler lock? I dea: use array of atomic b.

for flags

Avothel idea: have one A.B. to store "state" of lock A.B. locked: locked to three -7 to obtain set only obtain lock if - locked was false, and - I set it to t

Test and Set Lock

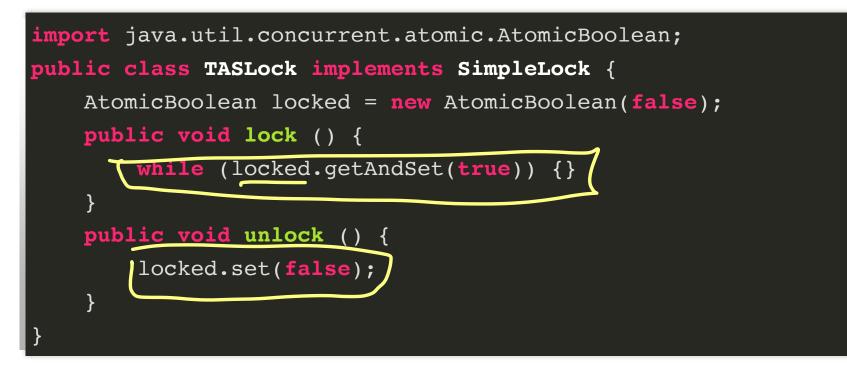
Idea. An AtomicBoolean locked stores state of the lock:

- locked.get() == true indicates the lock is in use
- locked.get() == false indicates the lock is free

Obtaining the lock:

- wait until locked is false, and set it to true Releasing the lock:
- set locked to false

```
TASLock in Code
```



download tas-locks.zip

Progress Guarantees

Question. Is TASLock deadlock-free? Starvation-free?

Alternative Implementation

Potential Issue:

- getAndSet operation is somewhat inefficient
 - slower than just get

Test and Test and Set Lock:

- check if locked
 - if not, attempt getAndSet
 - return if successful

TTASLock Implementation

```
public class TTASLock implements SimpleLock {
   AtomicBoolean locked = new AtomicBoolean(false);
   public void lock () {
      while (true) {
         while (locked.get()) {};
         if (!locked.getAndSet(true)) { return;}
      }
   }
   public void unlock() { locked.set(false);}
}
```

Comparing Efficiency

• tas-locks.zip